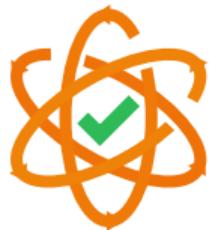


# Zoo : Un cadiciel pour la vérification de programmes OCaml 5 concurrents en logique de séparation

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30 janvier 2025

Vérification de programmes OCaml 5 *concurrents*.



Saturn  
Kcas



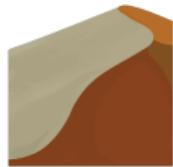
# À la recherche d'un langage de vérification

langage	concurrence	Iris	$\simeq$ OCaml	traduction	automatisation
Cameleer	:(	:(	:)	:)	:(
coq_of_ocaml	:(	:(	:)	:)	:(
CFML	:(	:(	:)	:)	:(
Osiris	:(	:(	:)	:)	:(
HeapLang	:(	:(	:(	:(	:(
Zoo	:(	:(	:)	:)	:(

## Zoo en pratique



OCaml



DUNE

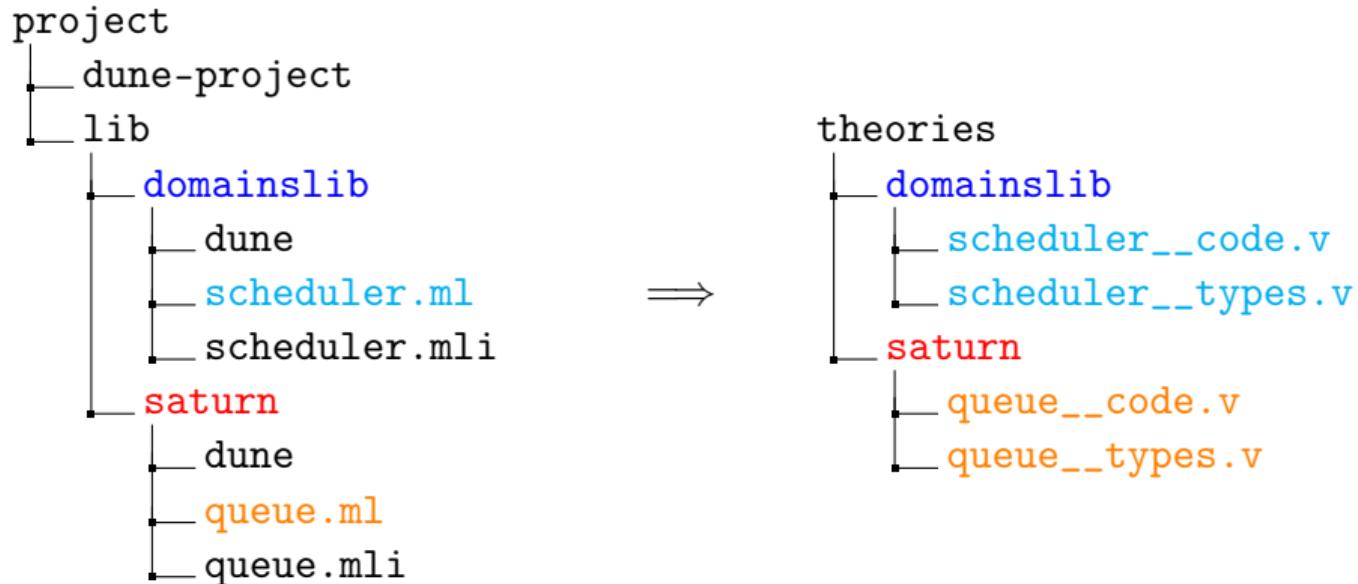
ocaml2zoo  
→



Zoo

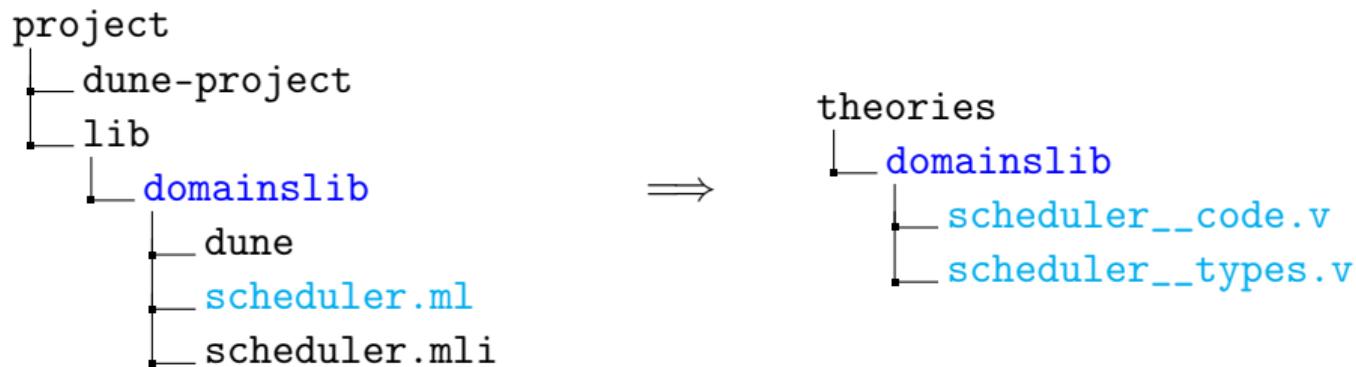


## Zoo en pratique



```
$ ocaml2zoo project theories
```

## Zoo en pratique



```
$ ocaml2zoo project theories
```

## Zoo en pratique

```
Lemma stack_push_spec_seq t  $\iota$  v :  
{{{  
  stack_model t vs  
}}} } }  
  stack_push t v  
{{{  
  RET ();  
  stack_model t (v :: vs)  
}}}.  
Proof.  
...  
Qed.
```

```
Lemma stack_push_spec_atomic t  $\iota$  v :  
<<<  
  stack_inv t  $\iota$   
|  $\forall$  vs,  
  stack_model t vs  
>>>  
  stack_push t v @  $\uparrow$  $\iota$   
<<<  
  stack_model t (v :: vs)  
| RET (); True  
>>>.  
Proof.  
...  
Qed.
```

# Types algébriques de données

```
type 'a t =
| Nil
| Cons of 'a * 'a t

let rec map fn t =
  match t with
  | Nil -> Nil
  | Cons (x, t) ->
    let y = fn x in
    Cons (y, map fn t)
```

```
Notation "'Nil'" := (
  in_type "t" 0
)(in custom zoo_tag).

Notation "'Cons'" := (
  in_type "t" 1
)(in custom zoo_tag).
```

```
Definition map : val :=
rec: "map" "fn" "t" =>
  match: "t" with
  | Nil => §Nil
  | Cons "x" "t" =>
    let: "y" := "fn" "x" in
    'Cons( "y", "map" "fn" "t" )
end.
```

## Enregistrements

```
type 'a t =
{ mutable f1: 'a;
  mutable f2: 'a;
}
```

```
let swap t =
  let f1 = t.f1 in
  t.f1 <- t.f2 ;
  t.f2 <- f1
```

```
Notation "'f1'" := (
  in_type "t" 0
)(in custom zoo_field).
Notation "'f2'" := (
  in_type "t" 1
)(in custom zoo_field).
```

```
Definition swap : val :=
  fun: "t" =>
    let: "f1" := "t".{f1} in
    "t" <-{f1} "t".{f2} ;;
    "t" <-{f2} "f1".
```

## Enregistrements anonymes

```
type 'a node =
| Null
| Node of
  { mutable next: 'a node;
    mutable data: 'a;
  }
```

```
Notation "'Null'" := (
  in_type "node" 0
)(in custom zoo_tag).
Notation "'Node'" := (
  in_type "node" 1
)(in custom zoo_tag).

Notation "'next'" := (
  in_type "node__Node" 0
)(in custom zoo_field).
Notation "'data'" := (
  in_type "node__Node" 1
)(in custom zoo_field).
```

## Fonctions mutuellement récursives

```
Definition f_g := (
  recs: "f" "x" => "g" "x"
  and:  "g" "x" => "f" "x"
)%zoo_recs.

(* boilerplate *)

let rec f x = g x
and g x = f x

Definition f := ValRecs 0 f_g.
Definition g := ValRecs 1 f_g.

Instance : AsValRecs' f 0 f_g [f;g].
Proof. done. Qed.
Instance : AsValRecs' g 1 f_g [f;g].
Proof. done. Qed.
```

# Concurrency

Atomic.set e <sub>1</sub> e <sub>2</sub>	e <sub>1</sub> <- e <sub>2</sub>
Atomic.exchange e <sub>1</sub> e <sub>2</sub>	Xchg e <sub>1</sub> .[contents] e <sub>2</sub>
Atomic.compare_and_set e <sub>1</sub> e <sub>2</sub> e <sub>3</sub>	CAS e <sub>1</sub> .[contents] e <sub>2</sub> e <sub>3</sub>
Atomic.fetch_and_add e <sub>1</sub> e <sub>2</sub>	FAA e <sub>1</sub> .[contents] e <sub>2</sub>
type t = { ...; mutable f: τ [@atomic]; ... }	
Atomic.Loc.exchange [%atomic.loc e <sub>1</sub> .f] e <sub>2</sub>	Xchg e <sub>1</sub> .[f] e <sub>2</sub>
Atomic.Loc.compare_and_set [%atomic.loc e <sub>1</sub> .f] e <sub>2</sub> e <sub>3</sub>	CAS e <sub>1</sub> .[f] e <sub>2</sub> e <sub>3</sub>
Atomic.Loc.fetch_and_add [%atomic.loc e <sub>1</sub> .f] e <sub>2</sub>	FAA e <sub>1</sub> .[f] e <sub>2</sub>

<https://github.com/ocaml/ocaml/pull/13404>  
<https://github.com/ocaml/ocaml/pull/13707>

## Bibliothèque standard

- ▶ Array
- ▶ Dynarray
- ▶ List
- ▶ Stack
- ▶ Queue
- ▶ Deque
- ▶ Domain
- ▶ Atomic\_array
- ▶ Mutex
- ▶ Condition

## Égalité physique : pile de Treiber

```
type 'a t =
  'a list Atomic.t

let create () =
  Atomic.make []

let rec push t v =
  let old = Atomic.get t in
  let new_ = v :: old in
  if not @@ Atomic.compare_and_set t old new_ then (
    Domain.cpu_relax ();
    push t v
  )
```

## Partage

```
let test1 = Some 0 == Some 0 (* true *)
let test2 = [0;1] == [0;1] (* true *)
```

## Conflits de représentation des valeurs

```
let test1 = Obj.repr false == Obj.repr 0 (* true *)
let test2 = Obj.repr None == Obj.repr 0 (* true *)
let test3 = Obj.repr [] == Obj.repr 0 (* true *)
```

## Partage + conflits

```
type any =
  Any : 'a -> any

let test1 = Any false == Any 0 (* true *)
let test2 = Any None == Any 0 (* true *)
let test3 = Any [] == Any 0 (* true *)
```

## Pile de Treiber

```
let rec push t v =
  let old = Atomic.get t in
  let new_ = v :: old in
  if not @@ Atomic.compare_and_set t old new_ then (
    Domain.cpu_relax ();
    push t v
  )
```

## Égalité physique : Eio.Rcfcd

```
type state = Open of Unix.file_descr | Closing of (unit -> unit)
type t = { mutable ops: int [@atomic]; mutable state: state [@atomic] }

let make fd = { ops= 0; state= Open fd }

let closed = Closing (fun () -> ())
let close t =
  match t.state with
  | Closing _ -> false
  | Open fd as prev ->
    let close () = Unix.close fd in
    let next = Closing close in
    if Atomic.Loc.compare_and_set [%atomic.loc t.state] prev next then
      ...
    else
      false
```

# Départage

```
let x = Some 0  
let test = x == x (* false *)
```



**Clément Allain**  
Impossible ! Identité unique.



**Armaël Guéneau**  
Ce serait du départage.



**Vincent Laviron**  
C'est possible !

## Eio.Rcfd

```
let closed = Closing (fun () -> ())
let close t =
  match t.state with
  | Closing _ -> false
  | Open fd as prev ->
    let close () = Unix.close fd in
    let next = Closing close in
    if Atomic.Loc.compare_and_set [%atomic.loc t.state] prev next then
      ...
    else
      false
```

## Constructeurs génératifs

```
type 'a liste =
| Nil
| Cons of 'a * 'a liste [@generative]

type state =
| Open of Unix.file_descr [@generative] [@zoo.reveal]
| Closing of (unit -> unit)
```

Merci de votre attention !